



Load-Time Security Certification for Real Smart-Cards

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Evaluation by B. Chetali, Q-H. Nguyen TrustedLabs/Gemalto (FR)

This talk



- How to design lightweight yet flexible and effective access control framework in a very restricted environment (Java Card)
- How to integrate the framework on a real card
- Bonus: demo of the prototype



Agenda



- Motivations and the Security-by-Contract idea
- The Java Card Background
- Contracts
- A (thin) hint of theory
- A (larger) taster of engineering
- Demo
- Conclusions



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Mobile payments



- NFC technology as enabler
- Secure element for storing secrets

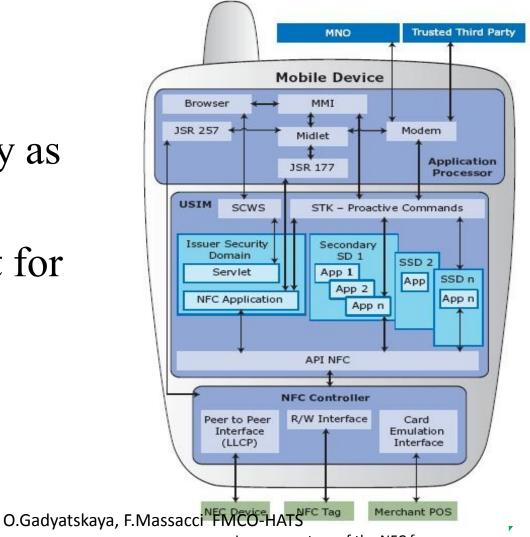


Image courtesy of the NFC forum

Pros of each secure element technology

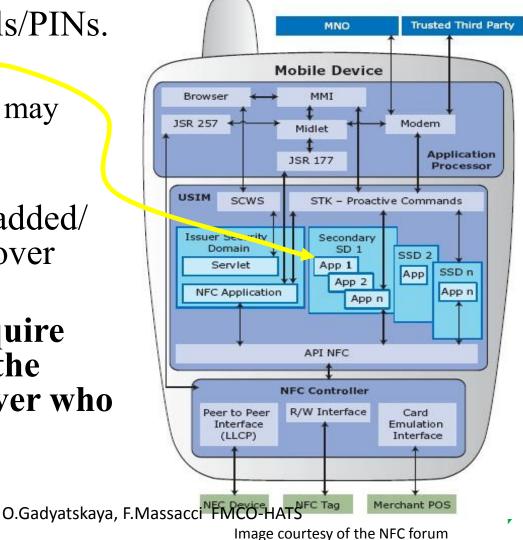




SIM as secure element



- Not only credentials/PINs. These are apps!
 - and some of them may even interact
 - this is Java
- New apps may be added/ old ones removed over time
- Sensitive apps require strict control (on the secure element) over who talks to whom



Design goals



• We need an on-card system that:

- Allows to add or remove applets
- Enables applets to declaratively control access to their shared resources (services)
- The access control policy can mention arbitrary applet identifiers (AIDs)
- The applet bytecode is validated by the card itself to respect the policies of other applets on card



Design constraints



• No modifications to the standard loading protocol, run-time environment or the virtual machine

Too expensive

• Most part of the trusted computing base is in ROM

- Cannot be modified after the card is in the field

- Applet providers can set up their policies independently
 - Telco does not want to be bothered



Existing solutions for Java Card:

- Can verify full information flow, but for predefined set of applets and off-card
- Can verify transitive control flow on card, but only for predefined and limited set of domains (applet owners)
- [Java Card protection] The policies are embedded into the applet code.



The threat model



- We assume an attacker that can:
 - Load or remove her applets on the card
 - Update access control policy of her own applets
- The attacker cannot:
 - Force loading or removal of someone else's applets or change their policies
 - Spoof someone else's applets pretending to be their legitimate owner
- The attacker's goal
 - Enable her applets to access illegally sensitive services of other applets



The Security-by-Contract idea



• SxC – particular instance of Load Time Verification

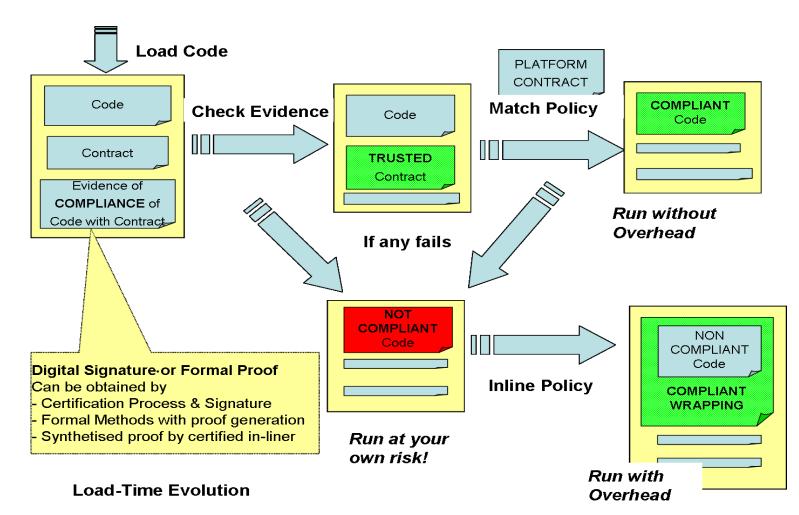
 Derived from Proof carrying code and Model carrying code ideas

• Well-tested for mobile platforms

- Java & .NET implementation (2008)
- Android (Manifest) implementation (Enck et al, 2010)
- But a smartphone isn't a card...

SxC workflow on mobile

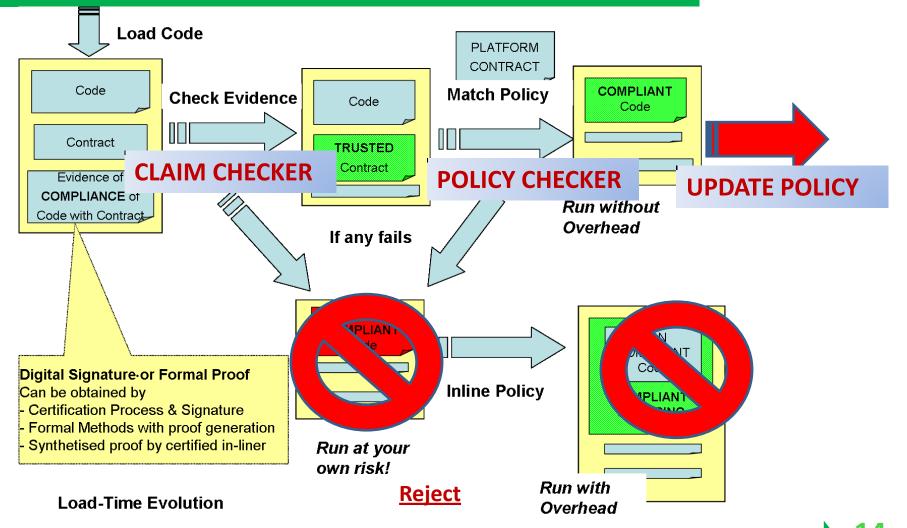




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S_xC workflow on smart cards





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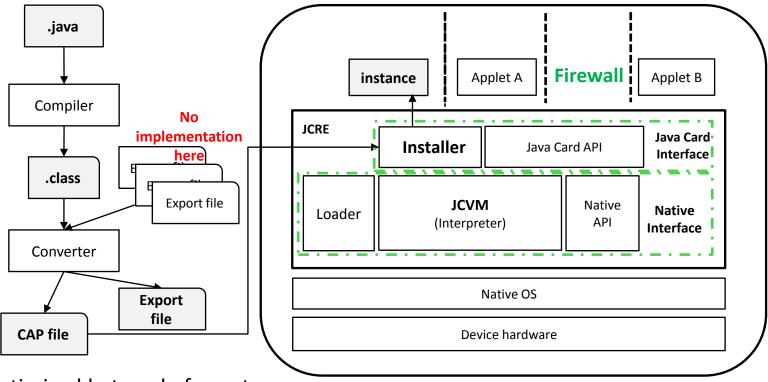


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The Java Card platform





Optimized bytecode format

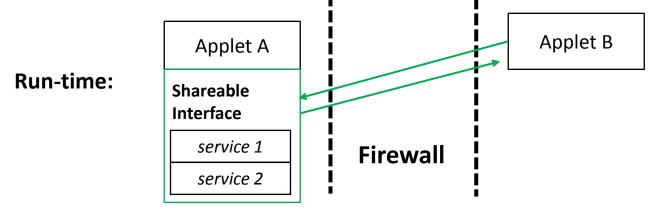
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How does JC really work?



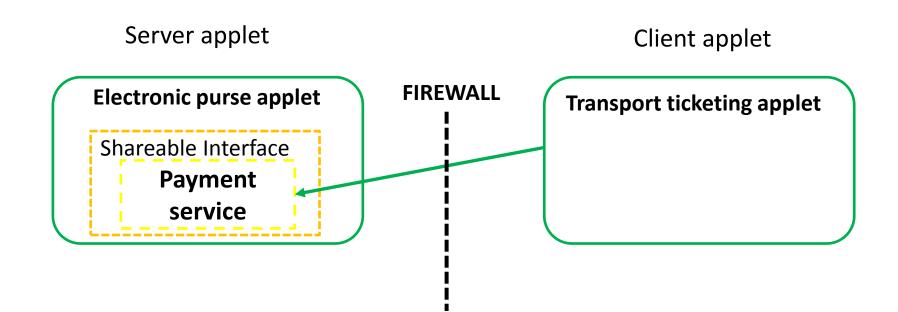


Access control is embedded into functional code

- Technical Consequence $1 \rightarrow$ If A checks who calls it, the access control policy cannot be updated unless the code is updated - sometimes code updates are not even possible
- Technical Consequence 2 \rightarrow If A does not check, then everybody can use it 04/10/2012

Example







ePurse applet: the ACL in the code



```
01 byte ClientsNumber = 1;
02 byte [] TransportAIDset =
\{0x01, 0x02, 0x03, 0x04, 0x05, 0x0C, 0x0A\};
03 final AID TransportAID = JCSystem.lookupAID
(TransportAIDset, (short)0, (byte) TransportAIDset.length);
04_ _
                                                  ACL checks
05 //the access control list
06 AID [] clientAIDs = {TransportAID};
07 //ACL check implementation
08 public short authorizedClient(AID clientAID) {
09
      for (short i=0; i<ClientsNumber; i++)</pre>
10
         if (clientAIDs[i].equals(clientAID))
11
            return i; //clientAIDs is in the ACL
12<sup>1</sup>
      return -1;
13<sup>}</sup>
```

ePurse applet: Shareable interface



14 //SI definition
<pre>15 public interface PaymentInterface extends Shareable {</pre>
<pre>16 //definition of the payment service</pre>
<pre>17 byte payment(short account_number); service</pre>
18 }
19 public class PaymentClass implements PaymentInterface {
<pre>20 byte payment_code = 0x08;</pre>
<pre>21 public byte payment(short account_number) {</pre>
<pre>22 //implementation of the service</pre>
AID clientAID = JCSystem.getPreviousContextAID();
if (authorizedClient(clientAID) == -1) //ACL check
return (byte) 0x00; //no service is provisioned
else return payment_code; //provision of the service
27 }
28 }
<pre>29 public PaymentClass PaymentObject;</pre>



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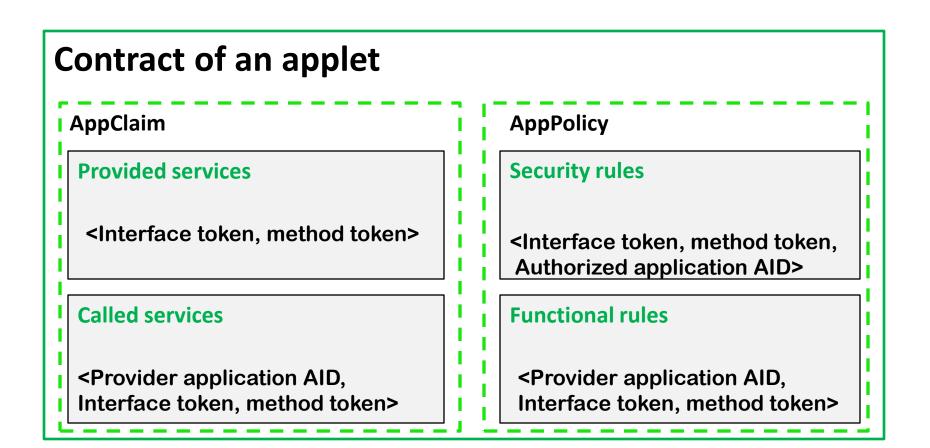
Contract I



- Applets come equipped with a contract
 Claims
 - I may **provide** these shareable interfaces with these services
 - I may **call** those methods from those interfaces
 - Security Rules
 - This service can only be called by this application
 - Functional Rules
 - I need these services from those applications
- When new applet arrives platform will check
 - contract complies with bytecode
 - contract is acceptable to other applets

Contract II





How do we get the tokens? - from Export files



Export file (snippet) of the Purse applet:

Service PaymentInterface.payment → gets token <0,0>



Invoked service tokens



Source code of Transport

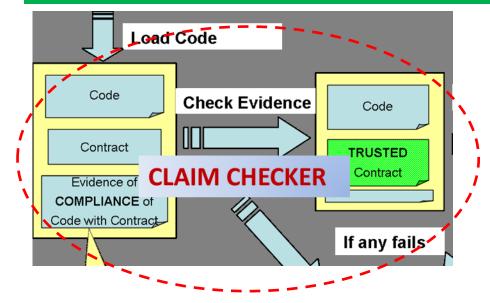
CAP file of *Transport*

<pre>01 private void connectServer(){ 02 final AID appletAID = JCSystem.lookupAID (serverAppletAID,(short)0,(byte)serverAppletAID .length);</pre>	<pre>package_info[1]{ AID_length 6 AID {1.2.3.4.5.b} }</pre>	Import component
<pre>03 if (appletAID == null) 04 ISOException.throwIt(ISO7816.SW_CONDITIONS_NOT_ SATISFIED); 05 PaymentObject = (PaymentInterface) (JCSystem.getAppletShareableInterfaceObject(app)</pre>	constantPool[16]{ external package_token 1 class_token 0	Constant Pool component
<pre>letAID, InterfaceDetails));</pre>	// bytecode of newE	Balance()
<pre>08 payment_code = PaymentObject.payment(account_number);</pre>	<pre>getstatic_a 17; getfield_b_this 2; invokeinterface 2 16 0; putfield_b 3; return;</pre>	Method component
09 return; 10 }	Called service <0,0> from AID 0x01020304050B	



The Claim Checker





Matches the Contract with the bytecode

For provided services:

➤ Checks the Shareable interfaces in CAP Export component For called services:

Finds all invokeinterface instructions (Method component and friends) and checks the invocation was declared

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Formally



- A deployed applet is a tuple <AID, Bytecode, ConstPool>
- A platform Θ is a set of currently deployed applets
- Security policy of the platform is a set of contracts $\{\{Contract_1\}, ..., \{Contract_N\}\}$ of currently deployed applets



Taxonomy of the JCVM instructions



Туре	Instructions
I	Arithmetic instructions; instructions that do not affect control flow. Cannot produce exceptions, execution proceeds to the next instruction: iadd
	Can throw run-time exceptions, but not security exceptions: irem
	Modify execution flow: goto, ifnull
IV	Return instructions: return
V	Can throw security exceptions: checkcast, iastore. The JCRE checks the object access rights here
VI	Invoke methods: invokeinterface, invokespecial, invokestatic, invokevirtual



The security theorem



- IF the JCRE is correct wrt specs:
 - [Firewall] applets only interact through Shareable interfaces
 - The Converter was correct and the CAP file was not tampered with
 - invokeinterface is the only invocation instruction that can be used for invoking services
- AND the SxC framework is correct wrt the specs
- THEN all methods invoked by any deployed applet B are authorized in the platform policy

Proof goes by cases of method invocation on the platform and inductively over the length of platform execution.



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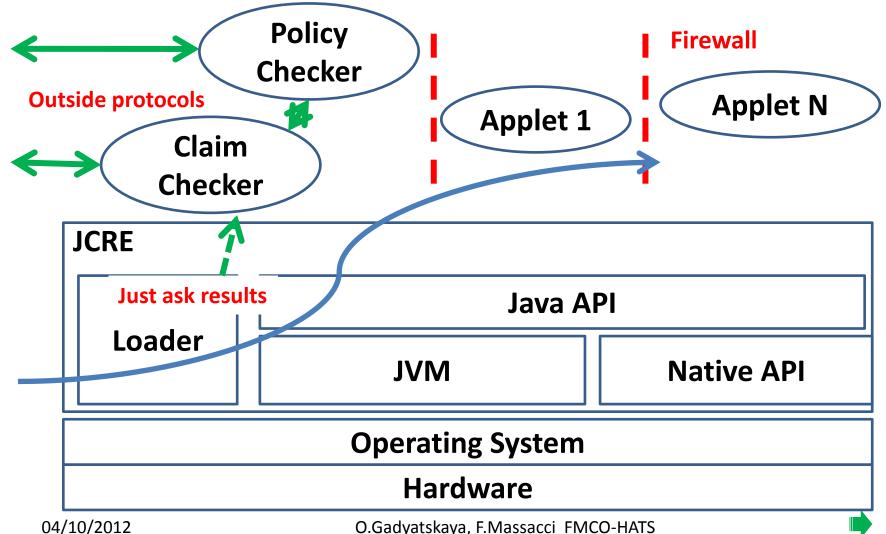


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Our first architecture: "as-onmobile"

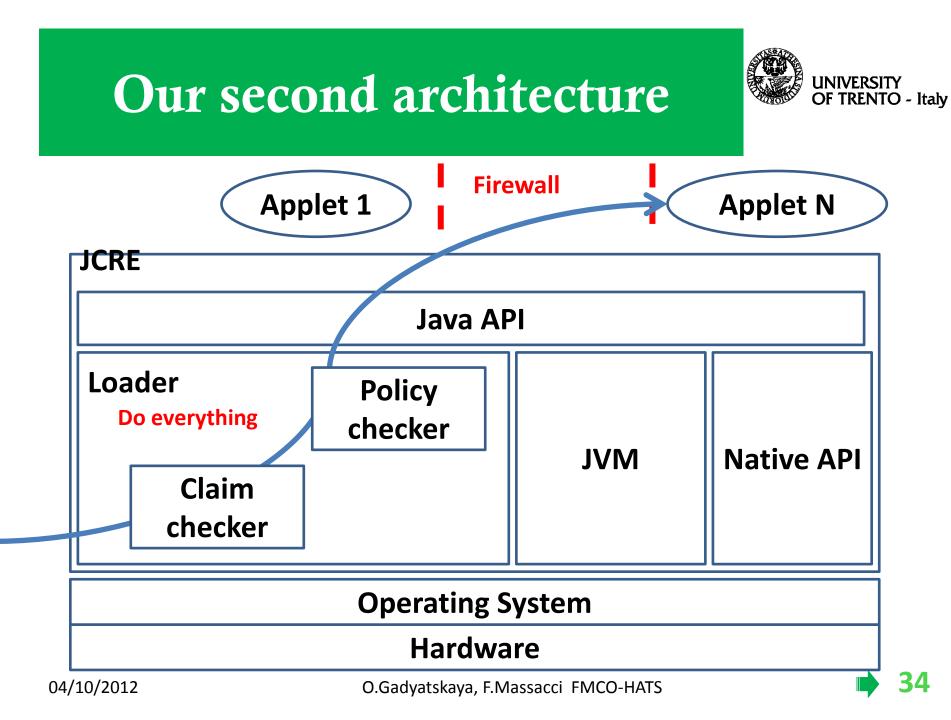




First engineering problem



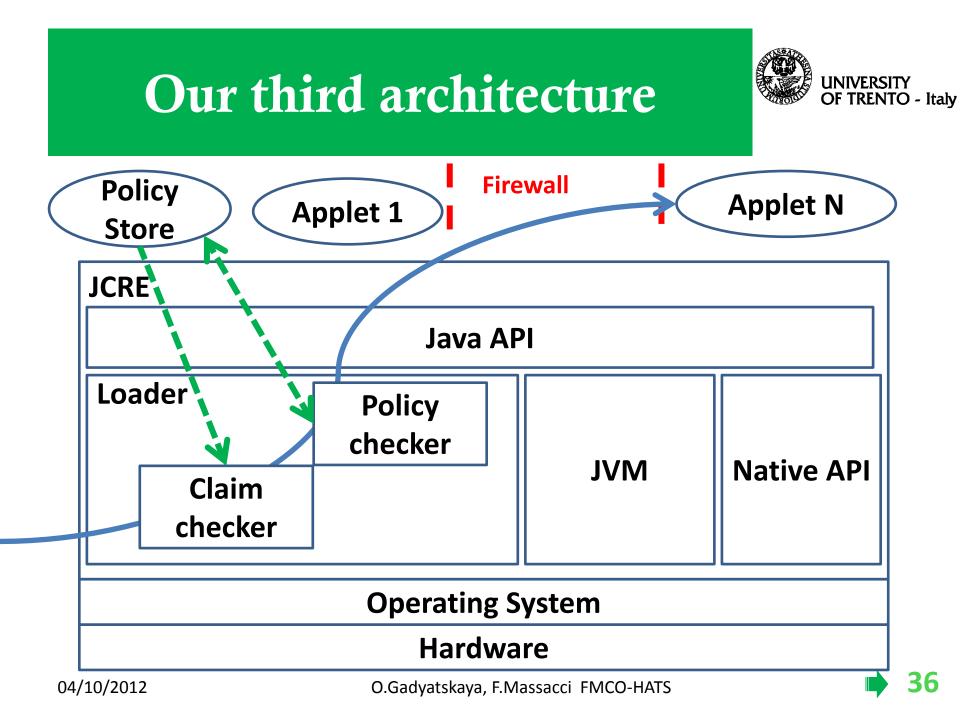
- We implemented Policy Checker as an applet – Footprint of checker 11KB and contracts 2KB
- BUT requires changing existing protocols!
 - Loading protocol standard plus check results of 1+2
 - New protocol with policy checker
 - New protocol with claim checker
- Loader can trust Policy Checker, but Claim Checker?
 - Needs signatures and certification
 - Too small improvement to justify new protocols



Second Engineering Problem



- More Effective and Efficient
 - Checkers no longer trust external checks of code
 - Eliminate check of signature!
 - Both checkers can be implemented in C
- But where do we put the policy?
 - We need to retrieve it and store it somewhere...
 - But the Loader is "printed"
 - We could have a "static int policy[]" but that's not going to work in the ROM



Third Engineering Problem



- How to deliver the Contract to the Checkers?
 - Can't change the loading protocol
- Both Checkers need applet AID...
 - AIDs are "big" → don't want to use them in the algorithms
 - AIDs only known at loading time → can't "print" them in Loader
- A bit of help from the platform
 - AID are mapped into Package ID (much shorter)
 - But still you have rules for AIDs not yet on board

Third Engineering Idea



- Each applet includes contract in CAP file Custom component
 - No need to send it separately
 - Arrives and leaves with applet
 - Updates identical to old code updates
 - Enables backward compatibility for cards and applets
- Checkers do not need trust anyone
 - Contract update would anyhow require code check
- PolicyStore references applet contract with PID
 - Mapping table from PID to AID
 - Checkers only get short matrix with loaded PIDs



Security policy on the card

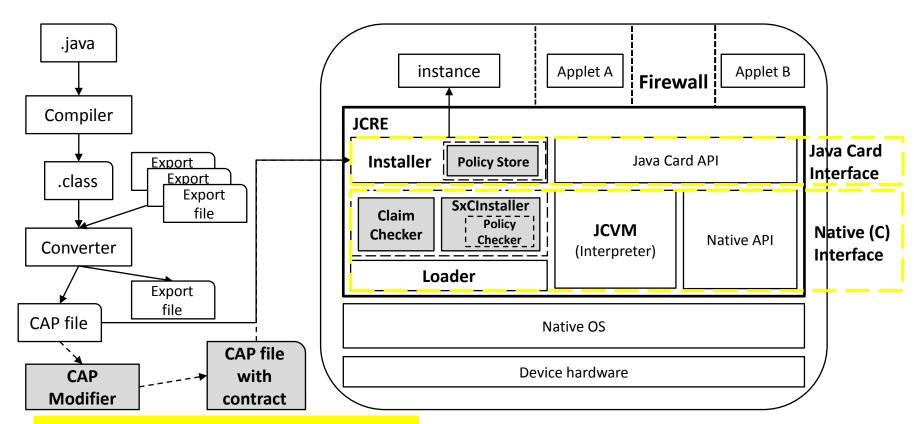


Arbitrary AIDs in the Mapping

	Policy on the card			
Small size and (frequent)	Policy (fixed size)	MayCall		
operations	All loaded contracts in an internal bit-arrays format	Possible future authorize for applets not yet on t		
		card	Big size and	
Big size and (rare) slow	Mapping Maintains correspondence between on-card IDs and AIDs	WishList Called services from ap	(rare) slow operations	
operations		not yet on the card		

The final architecture



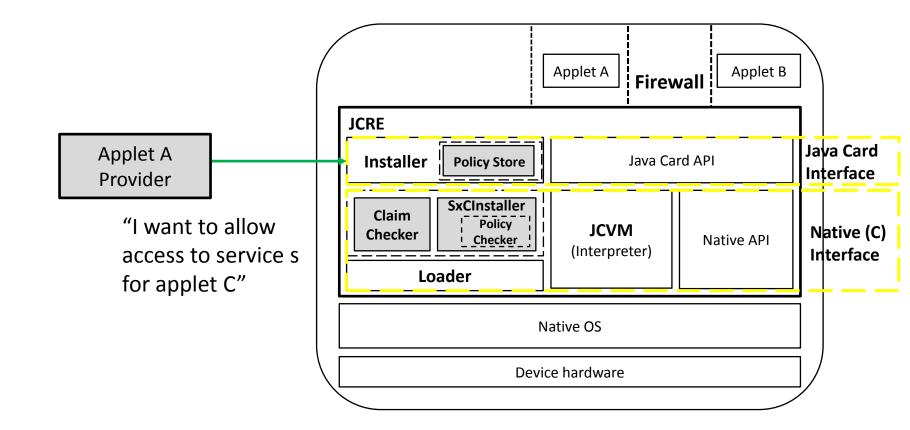


The SxC deployment process does not modify the standard Java Card tools



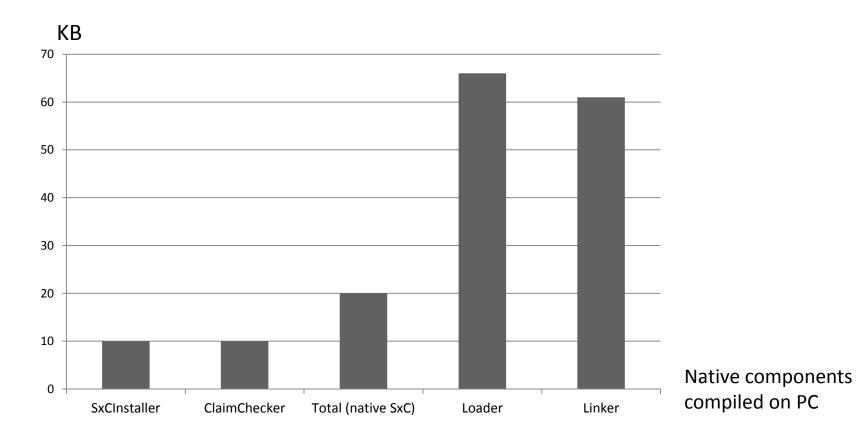
New applet policy update protocol





It is small enough





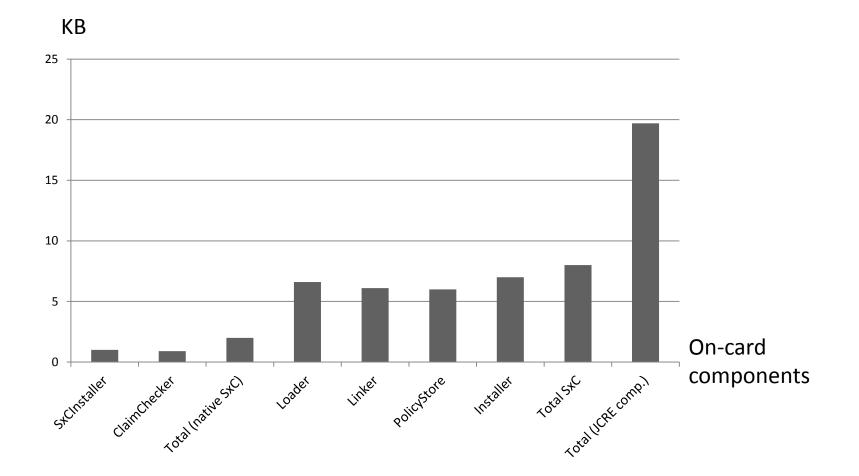
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It really works on the card





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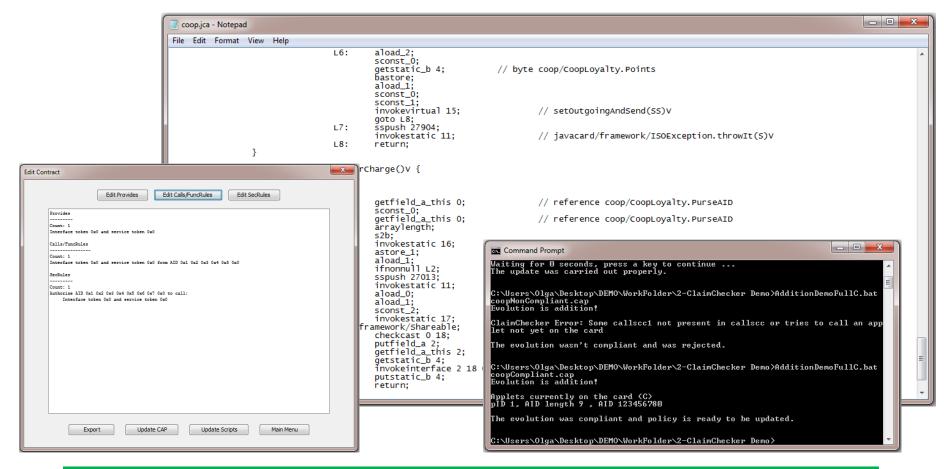
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Works on real applets



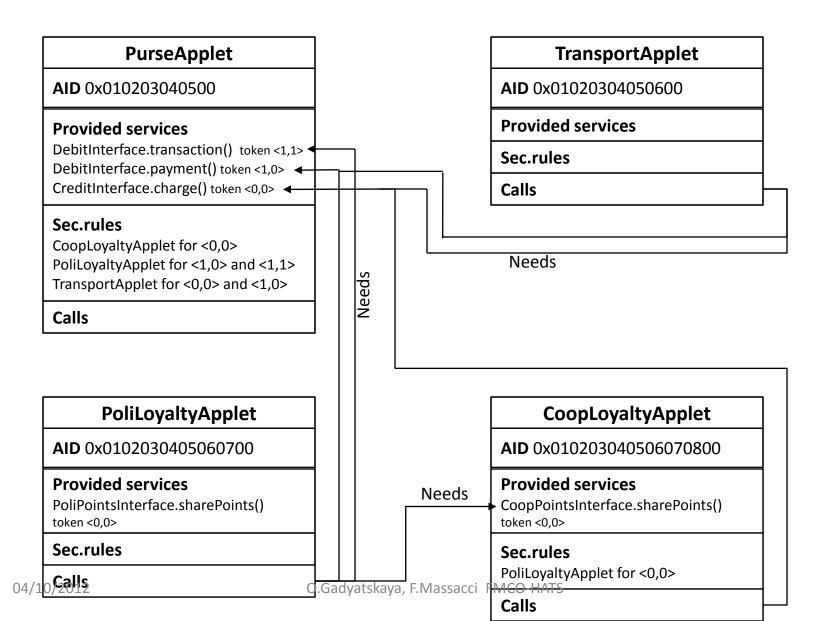
Applet	CAP file	# of methods	# of	LOCs
	\mathbf{size}	in CAP file	services	(.java)
Purse	2.5KB	6	1	66
Transport	$2.5 \mathrm{KB}$	5	0	92
EID	11.2KB	81	1	1419
ePurse	4.7KB	16	1	431







Demo scenario



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Industrial conclusions



- VISA is sceptical
- But
 - less sensitive applets require cheaper validation techniques



You can find more details in



- [POLICY'2011] N. Dragoni, E. Lostal, O. Gadyatskaya, F. Massacci, F. Paci: *A Load Time Policy Checker for Open Multi-application Smart Cards*
- [ICISS'2011] O. Gadyatskaya, E. Lostal, F. Massacci: Load Time Security Verification
- [BYTECODE'2012] O. Gadyatskaya, E. Lostal, F. Massacci: *Extended Abstract: Embeddable Security-by-Contract Verifier for Java Card*
- Some technical reports on my web page www.unitn.it/~gadyatskaya

Conclusions



- SxC framework performs loading time application certification
 - an applet is accepted only if it respects policies of other deployed applets
- Security code separated from the functional code
- It really works on a smart card
 - non-invasive addition to the standard Java Card deployment process







Send us your applets ...

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