



Autonomous and yet Secure Evolution for Smart Cards Applications

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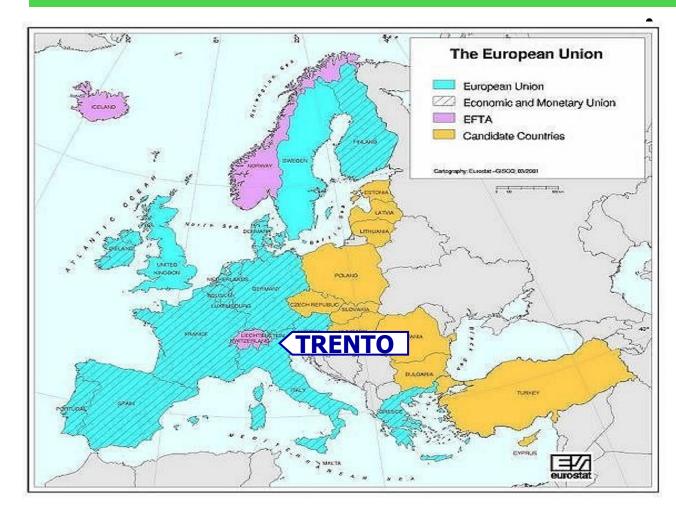
The talk plan



- Where's Trento?
- The rara avis of multi-application smart-cards
- Security-by-Contract for smart cards
- A (thin) slice of theory
- A (larger) slice of engineering
- Open problems

Trento in Space and Time





1962

 Institute of Social Science founded as locally funded Institution

1972

 Institute becomes private University

1982

 University becomes a state University with special autonomy

2001

 University becomes 1st in University Rankings

What do we do there?



- Organizational Level Security
 - Governance, Risk and Compliance (FM)
 - Security Requirements Engineering (FM,JM,PG)
- System Security
 - Run-time enforcement at ESB (FM,BC)
 - Browser Security (FM)
- Mobile/Embedded Code Security-by-Contract
 - Load-time security verification (FM)
 - Run-time information flow (BC)

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Smart cards today



- Modern computing devices
- Tamper-resistant security system
- Widely used
- But we have too many of them in our pockets



Open Multi-application smart cards



- Cards with multiple applets
 - allow post-issuance evolution (add/remove/update)
 - from different stakeholders
 - Asynchronously
- Interaction of applets on a single chip is natural:
 - Applications may interact exchanging loyalty points, transferring money or sharing valuable information.
- First paper I saw, I was a PhD Student 10yrs ago
 - Information Flow Verification for Multi-Applications Smart Card.
 - The Air-France, Hertz example...

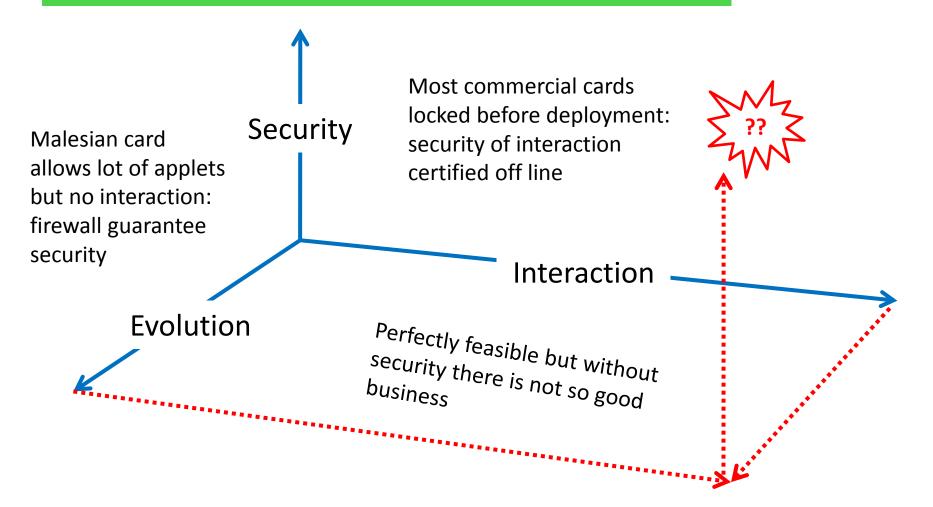
Java Card + GlobalPlatform



- GlobalPlatform = Middleware for secure management of applets (with open specification)
 - Lots of smart cards deployed with GP
- GlobalPlatform and JavaCard specifications
 - support loading, update and un-loading of many applications on the fly and asynchronously
 - allow interactions among applications (through services implementing Shareable interface)
- Still/Yet/But...
 - We don't really see multi-application cards in the wild.

What is there...





The usual evocative picture



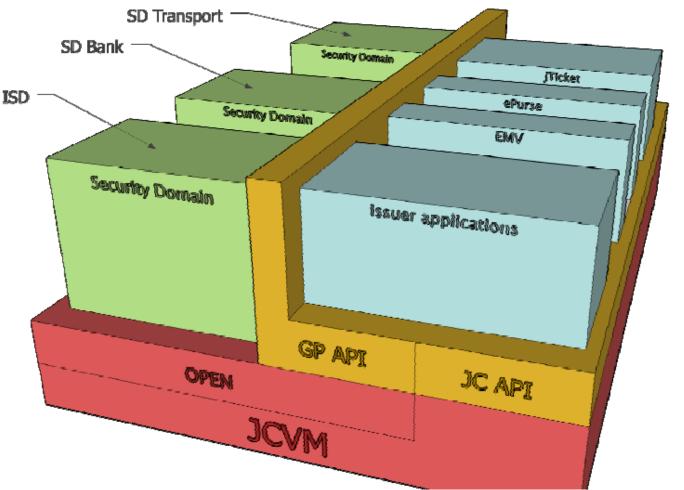
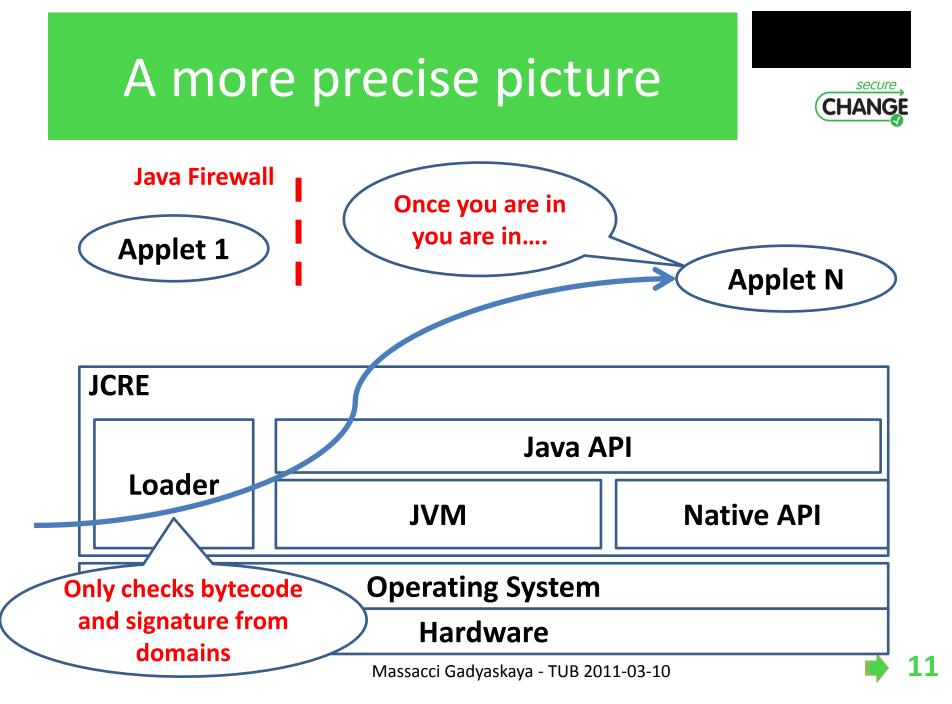


Image from D1.1 of SecureChange project

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How does Java Card firewall work?



- Applets interact through firewall using shareable interfaces
- Application ePurse of Bank
 - offers a service transfer_money.
 - does a preliminary access control checking caller AIDs in a list
- Application jTicket of Transport
 - wants to use transfer_money of ePurse
- What happens
 - jTicket asks the firewall for a reference to transfer_money.
 - Firewall passes call to ePurse. If jTicket is in the list, ePurse will return a reference to transfer_money service.
- Consequences
 - jTicket got a reference \rightarrow can use service from now on
 - ePurse wants to prevent jTicket use its service → must update itself

→ Business Model of Multi-Application SC (E+I+S) not supported

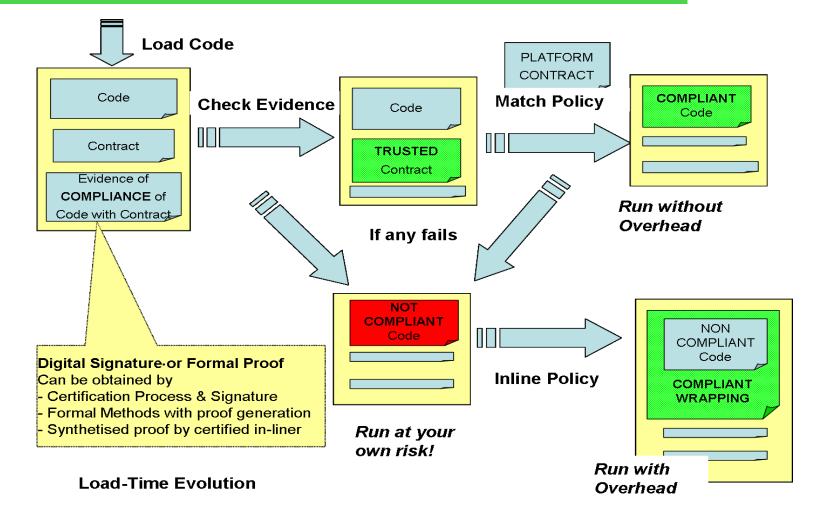
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Security-by-Contract idea

CHANGE



SxC as Load-time verification



- General idea of SxC for mobile devices:
 - Application has to be compliant with security policy of the device
 - Derived from PCC and MCC

Well-tested for mobile platforms

- Java & .NET implementation
- Eu S3MS project
- Many publication: JCS, JLAP, Comp. & Security, SCP, Elsevier IITR
- Policy checker could even run a small model checker
 - "allowed file.size > 1024Kb " vs "filesize < 512kb"
- But here we have a problem
 - Who sets the policy of device?
 - "Clear" for mobiles: operator, manfacturer, user

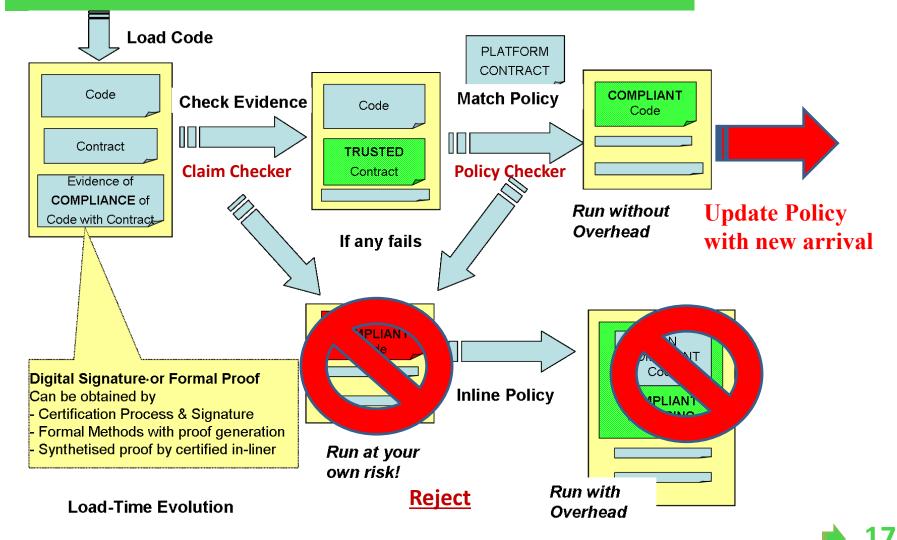
SxC for Smart Cards



- Whose policy?
 - The union of the policies of all applets
- Broader Contract
 - Claims
 - I may provide these shareable interfaces
 - I may call those methods from those interfaces
 - Security Rules
 - I can only be called by this Application/Package
 - Functional Rules
 - I need these methods from those interfaces

S_xC workflow for smart cards





SxC Example



 Already installed Applet ePurse with Contract:

- Applet jTicket arrives with Contract:
- jTicket is loaded, cheked, and finally installed.

- Applet i-Travel arrives with Contract:
- i-Travel is rejected: load process is not committed

- Provides = {*transferMoney*}
- Calls = {}
- Sec.rules = {transferMoney -> {jTicket}}
- Func.rules = {}
 - Provides = {ageDiscount, loyaltyDiscount}
 - Calls = {ePurse.transferMoney}
 - Sec.rules = { ageDiscount → {IDapplet}, loyaltyDscount → {ePurse}}
 - Func.rules = {ePurse.transferMoney}
- Provides = {}
- Calls = {ePurse.transferMoney}
- Sec.rules = {}
- Func.rules ={}

Formal Model of a JC Platform



Platform Θ =

- <Δ_A, Δ_S, *A*, shareable(), invoke(), sec.rules(), func.rules()>
 - $-\Delta_A =$ domain of applications, $\Delta_S =$ domain of services $-\mathcal{A} \subseteq \Delta_A$
 - applets deployed (installed) on the platform
 - shareable(), invoke(): $\Delta_A \rightarrow p(\Delta_S)$
 - Services offered by applet (resp. invoked by applet)
 - sec.rules(): $\Delta_A \times \Delta_S \rightarrow p(\Delta_A)$
 - For any applet and its services which applets can call it
 - func.rules(): $\Delta_A \rightarrow p(\Delta_S)$
 - Services that must be present in order for the applet to function

Why we use different names?



• Platform has

– Shareable(A) $\subseteq \Delta_s$ and invoke(A) $\subseteq \Delta_s$

Contract has

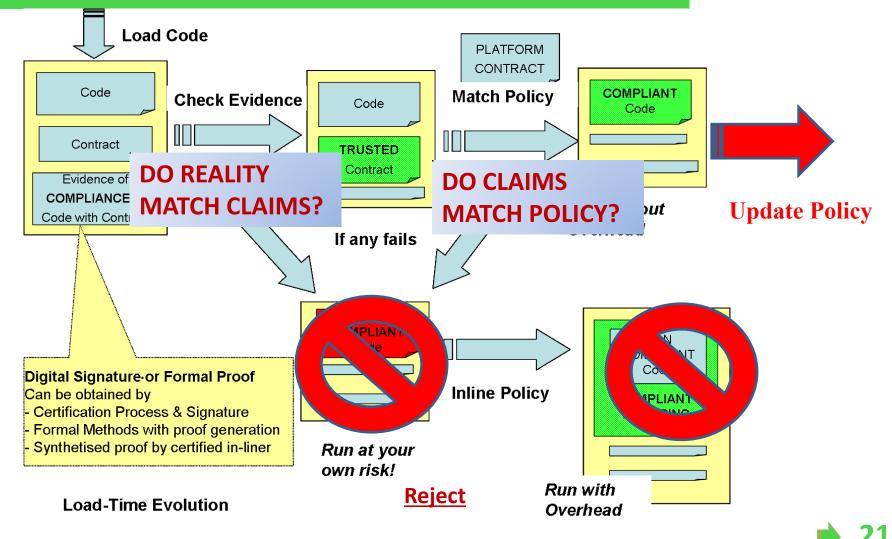
– Provides(A) $\subseteq \Delta_s$ and Calls(A) $\subseteq \Delta_s$

- Same difference between reality and claims
 - The first is reality, what really is there
 - The seconds are the claims, they might be honest but might also not correspond to truth



S_xC workflow for smart cards

CHANGE



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Introducing evolution to the model



- Let B be an application, an evolved platform Θ' for B from a platform Θ is defined according to the next types of changes:
 - B is a new applet to be added to the platform,
 - old applet B is removed from the platform,
 - update of an installed applet B
 - Add/remove of a service to shareable(B)
 - Add/remove of a service to invoke(B)
 - Add/remove of an access authorization to sec.rules(B)
 - Add/remove of a service to func.rules(B)

Checking Changes Incrementally



- For each type of change the Claim Checker and the Policy checker should verify only the parts of the platform that are touched by changes.
- For new applet B:
 - Claim Checker has to verify that
 - shareable(B)=ProvidesB
 - invoke(B)=CallsB
 - (or to extract shareable(B) and invoke(B) from the code and write these sets into the ContractB)
 - The Policy Checker has to check that for all applets $A \in A$:
 - if $A.s \in CallsB$ then $(s,B) \in sec.rules(A)$
 - if $A.s \in func.rules(B)$ then $s \in ProvidesA$
 - if B.s \in CallsA then (s,A) \in sec.rules(B)

Trickier Example



• Applet ePurse:

• Applet jTicket:

Provides = {transferMoney}
Calls = {}
Sec.rules = {transferMoney → {jTicket}}
Func.rules = {}

- Provides = {ageDiscount, loyaltyDiscount}
- Calls = {*ePurse.transfer_money*}
- Sec.rules ={ ageDiscount → {IDapplet}, loyaltyDiscount → {ePurse}}
- Func.rules = {ePurse.transfer_money}
- Now we update ePurse
- What happens?

- Provides = {*transferMoney*}
- Calls = { jTicket.ageDiscount }
- Sec.rules = {transferMoney → {jTicket}}
- Func.rules = {}

Secure Platform



- A platform O remains secure during evolution
 - This is what you really want after each update
 - For every applet the traces of real executions respects its security and functional rules
 - Whenever somebody calls you it is authorized
 - Whenever you need to call an essential service it is still there (provided it was there before)
- Security and functionality in terms of Contracts
 - Contracts do not violate Global Policy
 - Claims are consistent with bytecode
 - Otherwise update is rejected
- Need to show the two coincide.

Security Theorem



- IF Platform was secure before the update,
- & IF shareable interfaces are only means for inter-app communication
- & IF Claim Checker and the Policy Checker are sound and accepted an update at the loading time,
- THEN evolved platform will be secure.
 - Proving by contradiction that if security or functionality is broken on the platform, then either the ClaimChecker, or the Policy Checker will reject the update
- Still an Engineering gap
 - In theory it could work for Application IDs in contracts,
 - in practice we may need to weaken the claim to Package Ids
- Depends on what we can implement in the claim checker

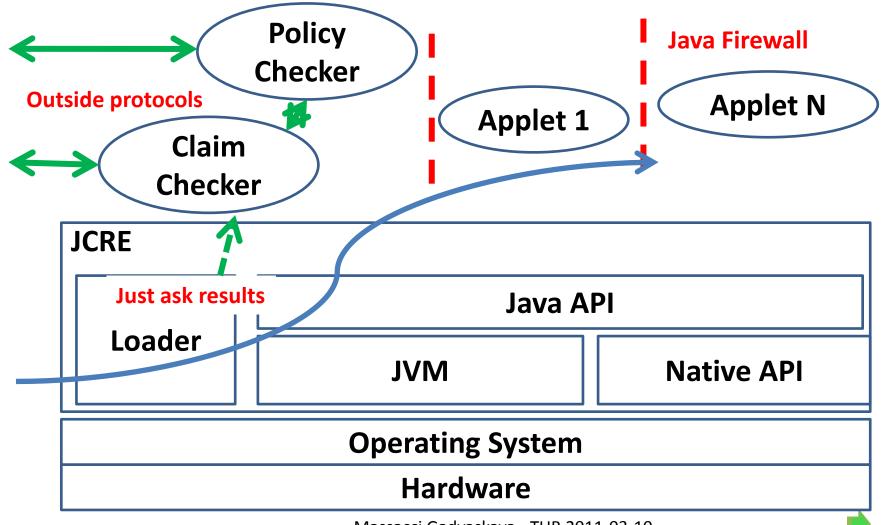
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Our First Architecture



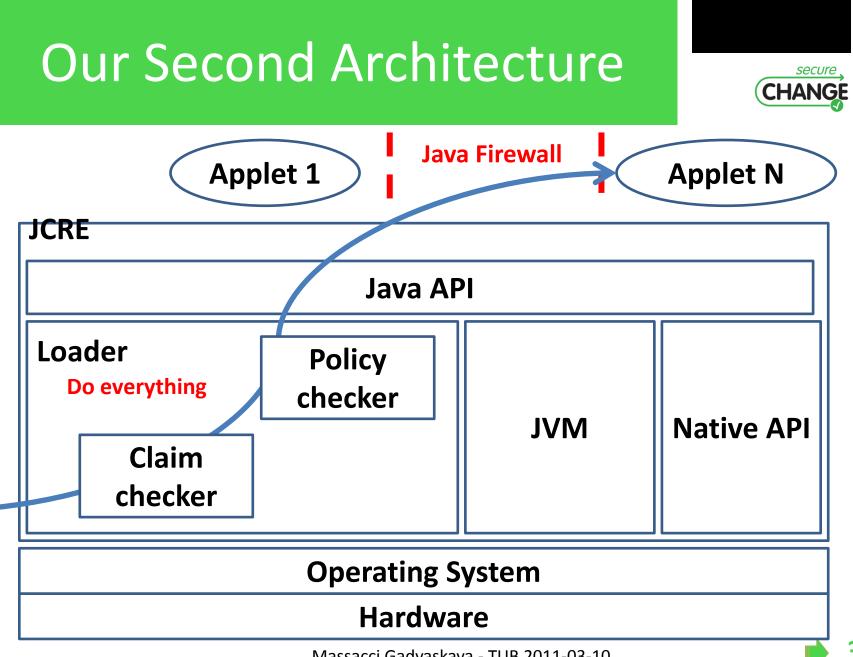


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First Engineering problem



- Implemented Policy Checker
 - POLICY'11 short paper
 - Footprint of checker 11KB and contracts 2KB
- Require changing existing update protocols
 - 1stprotocol with policy checker
 - 2nd protocol with claim checker
 - 3rd protocol is standard loading plus check results of 1+2
- Loader can trust policy checker, what about claim checker?
 - Needs signatures and certification
 - Too small improvement to justify change update protocol

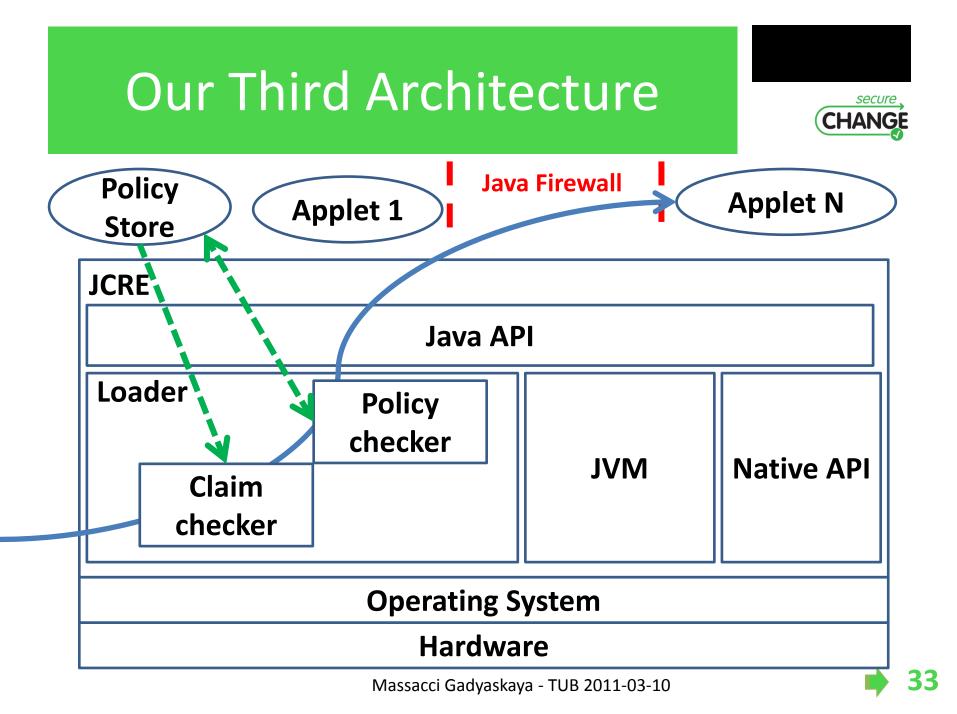


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Second Engineering Problem



- More Effective and Efficient
 - Loader no longer trust external checks of code
 - Eliminate checks of signatures beside standards
 - Both checkers can be implemented in C
- But where do we put the policy?
 - We need to retrieve it and store it somewhere...
 - but loader is NOT loaded in the EEPROM
 - We could have a "static int policy[]" but that's not going to work in the ROM

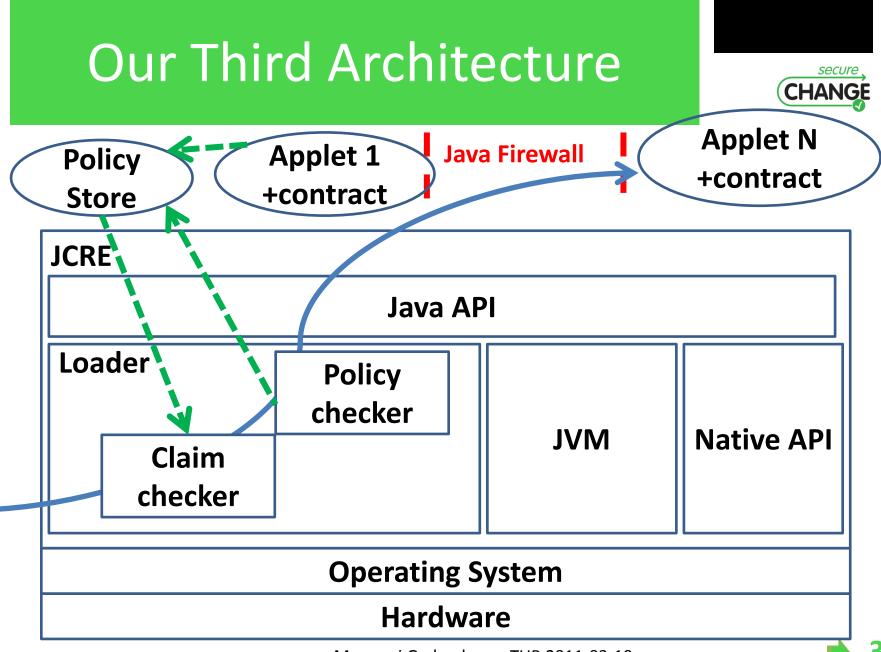


Third Engineering Problem



- C and Java don't mix well
 - The loader can "easily" invoke the Policy Store applet at the beginning of the process andpass reference to it to the loader
 - Just need a Java shell onto the loader
 - but how to tell it the result at the end??
 - It must be the checked contract and nothing else
- Who's giving the contract to the checker?

– Must change the protocol of update...



Engineering Idea



- Each Applet includes contract in java package
 - No need to send it separately
 - Arrives and leaves with applet
 - Neutral: contract update requires re-running claim checker
 - Cons: contract update requires code update
 - But in this way claim checker re-run is automatic!
- Policy store references applet contract
 - Keep efficiency of C implementation with Java flexibility
- Checkers do not need trust anyone
- Next validation by Smart card manufacturer

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Trickier Example



• Applet ePurse:

• Applet jTicket:

- We update ePurse
- Update is accepted
- What happens later if jTickets wants to drop access to ePurse?

Provides = {transferMoney}
Calls = {}
Sec.rules = {transferMoney → {jTicket}}
Func.rules = {}

- Provides = {ageDiscount, loyaltyDiscount}
- Calls = {*ePurse.transferMoney*}
- Sec.rules ={ ageDiscount → {IDapplet}, loyaltyDiscount → {ePurse}}
- Func.rules = {ePurse.transfer_money}
 - Provides = {*transferMoney*}
 - Calls = {jTicket.loyaltyDiscount}
 - Sec.rules = {transferMoney → {jTicket}}
 - Func.rules = {jTicket.loyaltyDiscount}

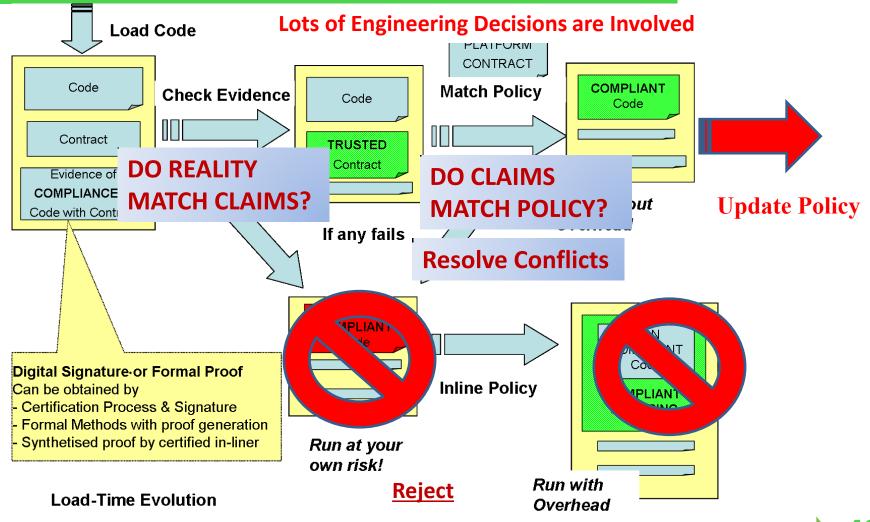
A Conflict Resolution Componnet?



- What happens if ePurse owner wants it to be removed from the platform?
 - jTicket needs the service ePurse.transfer_money
 - But ePurse doesn't want (now) to give him this
- Two possibilities:
 - to forbid ePurse to be removed OR
 - to remove ePurse and make jTicket unselectable.
- (Automatic) Conflict resolution requires investigation of stakeholders (security domains) hierarchy.

Conclusions: SxC for Smart-Cards









Send us your applets!

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